Robust mean change point testing in high-dimensional data with heavy tails

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Collaborators



Reference

Li, M.*, Chen, Y.*, Wang, T. and Yu, Y. (2023) Robust mean change point testing in high-dimensional data with heavy tails. *arXiv preprint*, arXiv: 2305.18987.

Consider the model



Entries of E are independent random variables with mean 0, variance 1 and distribution $P_e\!.$

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Task: mean change point testing

 H_0 : no change in the columns of θ vs. $H_1: \exists \text{ a change in the columns of } \theta$

Change point detection – testing problem

Model: $X = \theta + E \in \mathbb{R}^{p \times n}$

• Null hypothesis H_0 (no change)

$$H_0: heta \in \Theta_0(p,n) := ig\{ heta: heta_1 = heta_2 = \ldots = heta_n = \mu \in \mathbb{R}^p ext{ for some } \muig\}.$$

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• Alternative hypothesis H_1 (\exists change)

$$H_1: \theta \in \Theta(p, n, s, \rho) := \bigcup_{t_0=1}^{n-1} \Theta^{(t_0)}(p, n, s, \rho)$$

where

$$\begin{split} \Theta^{(t_0)}(p,n,s,\rho) &:= \Big\{ \theta : \theta_t = \mu_1 \text{ for } 1 \le t \le t_0, \ \theta_t = \mu_2 \text{ for } t_0 + 1 \le t \le n, \\ \underbrace{\|\mu_1 - \mu_2\|_0}_{\text{sparsity level}} \le s, \underbrace{\frac{t_0(n-t_0)}{n} \|\mu_1 - \mu_2\|_2^2}_{\text{normalised signal strength}} \ge \rho^2 \Big\}. \end{split}$$

Minimax testing rate

Definition. Let Φ be the set of all test functions $\phi : \mathbb{R}^{p \times n} \to \{0, 1\}$. Denote the minimax testing error

$$\mathcal{R}_{\mathcal{Q}}(\rho) := \inf_{\phi \in \Phi} \mathcal{R}_{\mathcal{Q}}(\rho, \phi)$$

$$:= \inf_{\phi \in \Phi} \left\{ \underbrace{\sup_{P_e \in \mathcal{Q}} \sup_{\theta \in \Theta_0(p,n)} \mathbb{E}_{\phi}}_{\mathsf{Type \, l \, error}} \underbrace{\mathbb{E}_{P_e \in \mathcal{Q}} \sup_{\theta \in \Theta(p,n,s,\rho)} \mathbb{E}(1-\phi)}_{\mathsf{Type \, l \, error}} \right\}.$$

$$v_{\mathcal{Q}}^*(p, n, s) \text{ is the minimax testing rate if}$$

$$1. \exists \text{ test } \phi, \text{ s.t. } \mathcal{R}_{\mathcal{Q}}(\rho, \phi) \leq 1/2 \text{ when } \rho^2 \gtrsim v_{\mathcal{Q}}^*(p, n, s),$$

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2. \forall test ϕ , have $\mathcal{R}_{\mathcal{Q}}(\rho, \phi) > 1/2$ when $\rho^{2} \gtrsim v_{\mathcal{Q}}(p, n, s)$.

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$$\begin{split} \mathcal{R}_{\mathcal{Q}}(\rho) &:= \inf_{\phi \in \Phi} \mathcal{R}_{\mathcal{Q}}(\rho, \phi) \\ &:= \inf_{\phi \in \Phi} \bigg\{ \underbrace{\sup_{P_e \in \mathcal{Q}} \sup_{\theta \in \Theta_0(p,n)} \mathbb{E}\phi}_{\text{Type I error}} + \underbrace{\sup_{P_e \in \mathcal{Q}} \sup_{\theta \in \Theta(p,n,s,\rho)} \mathbb{E}(1-\phi)}_{\text{Type II error}} \bigg\}. \\ v_{\mathcal{Q}}^*(p,n,s) \text{ is the minimax testing rate if} \end{split}$$

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$$\exists \text{ test } \phi, \text{ s.t. } \mathcal{R}_{\mathcal{Q}}(\rho, \phi) \leq 1/2 \text{ when } \rho^2 \gtrsim v_{\mathcal{Q}}^*(p, n, s),$$

- 2. \forall test ϕ , have $\mathcal{R}_{\mathcal{Q}}(\rho, \phi) > 1/2$ when $\rho^2 \lesssim v_{\mathcal{Q}}^*(p, n, s)$.
- The distribution of each entry in E belongs to some class Q.
- Liu et al. (2021) derived the minimax testing rate for $Q = \{N(0,1)\}$.
- Heavy-tailed distributions in Q?

Two types of heavy-tailedness:

Definition ($\mathcal{G}_{\alpha,K}$ class). For any $P \in \mathcal{G}_{\alpha,K}$ and r.v. $W \sim P$, $\mathbb{E}W = 0$, $\mathbb{E}W^2 = 1$ and $\mathbb{E}\exp\{|W/K|^{\alpha}\} \leq 2$.

Sub-Weibull distributions of order α ; possessing exponentially-decaying tails

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Distributions with **finite** α -**th moment**; possessing polynomially-decaying tails



$s_{\mathcal{P}}^* = p^{\frac{1}{2}} p^{-\gamma(\alpha)}$

 $s_{\mathcal{G}}^* = p^{\frac{1}{2}} \log^{-\beta(\alpha)}(ep)$

- Minimax rate upper bound v^{U} : construct a test procedure ϕ , such that $\mathcal{R}_{\mathcal{Q}}(\rho, \phi) \leq 1/2$ when $\rho^2 \geq v^{U}$.
- Minimax rate lower bound v^{L} : usually via Le Cam's two point method.

		Upper bound	Lower bound
$\mathcal{G}_{lpha,K}^{\otimes}$	Dense	(i) $\sqrt{p \log \log(8n)} + \log \log(8n)$	(ii) $\sqrt{p(\log\log(8n))^{\omega_1}} + \log\log(8n)$
	Sparse	(iii) $s \log^{2/\alpha}(ep/s) + \log \log(8n)$	(iv) $s \log^{2/\alpha}(ep/s) + \log \log(8n)$
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Note:
$$\omega_1 = \mathbb{1}_{\left\{s > \sqrt{p \log \log(8n)}\right\}}$$
 and $\omega_2 = \mathbb{1}_{\left\{s > \sqrt{p \log \log(8n)} \text{ and } \alpha \ge 4\right\}}$

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Matching rates!

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A minimax gap of $\sqrt{\log \log n}$

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Sub-Weibull $\mathcal{G}_{\alpha,K}$ – dense

Consider a dyadic grid $\mathcal{T} := \{1, 2, 4, \dots, 2^{\lfloor \log_2(n/2) \rfloor}\}$ and CUSUM-type statistics

$$Y_t := \frac{\sum_{i=1}^t X_i - \sum_{i=1}^t X_{n+1-i}}{\sqrt{2t}} \in \mathbb{R}^p.$$

Aggregation across coordinates:

$$A_t := \sum_{j=1}^p (Y_t^2(j) - 1).$$

Test:

$$\phi_{\mathcal{G},\text{dense}} := \mathbb{1}_{\{\max_{t \in \mathcal{T}} A_t > r\}}.$$

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Thresholding step

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Thresholding step + sample splitting

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Thresholding step + sample splitting

Test:

$$\phi_{\mathcal{G},\text{sparse}} := \mathbb{1}_{\{\max_{t \in \mathcal{T}} A_{t,a} > r\}}.$$

► $\mathcal{R}_{\mathcal{G}}(\rho, \phi_{\mathcal{G}, \text{sparse}}) \leq 1/2 \text{ as long as } \rho^2 \gtrsim s \log^{2/\alpha}(ep/s) + \log \log(8n).$

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In $\mathcal{P}_{\alpha,K}$, each entry of the noise matrix E has finite α -th moment. For high-dimensional mean change point testing problem:

- 1. When $\alpha \leq 4$, the sparse regime disappears.
- 2. Median-of-means-type statistics are effective in handling heavy-tailed data.
- 3. When $\alpha \ge 4$, in the sparse regime, we propose a computationally efficient test that achieves minimax optimality.

Transition boundary under $\mathcal{P}_{\alpha,K}$

Finite a-th moment 0.5 0.4 Sparse regime $\gamma(lpha) = rac{1}{lpha-2} \wedge rac{1}{2}$ 0.3 γ 0.2 -Dense regime 0.1 0.0 -2 8 α

 $s_{\mathcal{P}}^* = p^{\frac{1}{2}} p^{-\gamma(\alpha)}$

Consider the dense regime

$$s \ge p^{\frac{1}{2} - \left(\frac{1}{\alpha - 2} \wedge \frac{1}{2}\right)}.$$

For $\alpha \geq 2$, we show that $\mathcal{R}_{\mathcal{P}}(\rho) \geq 1/2$ whenever

$$\rho^2 \lesssim p^{(2/\alpha)\vee(1/2)} (\log\log(8n))^{\omega/2} + \log\log(8n))^{\omega/2}$$

with
$$\omega = \mathbb{1}_{\left\{s > \sqrt{p \log \log(8n)}\right\} \cap \{\alpha \ge 4\}}$$
.

• When $\alpha \leq 4$, the dense regime becomes $s \gtrsim 1$, i.e. no sparse regime.

Detour: Median-of-means (MoM)

Let $X_1, \ldots, X_n \in \mathbb{R}$ be i.i.d random variables with mean μ and variance σ^2 and consider the following **median-of-means estimator**

$$\hat{\mu}_{\text{MoM}} = \text{median}\left(\frac{1}{m}\sum_{i=1}^{m}X_i, \dots, \frac{1}{m}\sum_{i=(k-1)m+1}^{km}X_i\right).$$

Let $\delta \in (0,1), k = \lceil 8 \log(1/\delta) \rceil$ and m = n/k. Then w.p. at least $1 - \delta$,

$$|\hat{\mu}_{\mathrm{MoM}} - \mu| \le \sigma \sqrt{\frac{32\log(1/\delta)}{n}}.$$

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• δ is an input to the estimator, through k (number of groups).

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- 'sub-Gaussian' property.
- δ is an input to the estimator, through k (number of groups).
- For a given δ , the result is only possible when n is at least $8 \log(1/\delta)$.
- Equivalently, for *n* fixed, δ needs to be larger than $\exp(-n/8)$.

Finite moment $\mathcal{P}_{\alpha,K}$ – **dense**

• (In sub-Weibull dense...) For $t \in \mathcal{T}$, we compute

$$Y_t := \frac{\sum_{i=1}^t X_i - \sum_{i=1}^t X_{n+1-i}}{\sqrt{2t}} \in \mathbb{R}^p.$$

and

$$A_t := \sum_{j=1}^p (Y_t^2(j) - 1).$$

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$$Z_i := (X_i - X_{n-i+1})/\sqrt{2}.$$

For $t \in \mathcal{T}$, split $\{Z_1, \ldots, Z_t\}$ into G_t groups of equal size

$$\mathcal{Z}_{t,1}, \mathcal{Z}_{t,2}, \ldots, \mathcal{Z}_{t,G_t}.$$

Each group contains t/G_t elements.

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Each group contains t/G_t elements.

• Set
$$V_{t,g} \in \mathbb{R}^p$$
 with $V_{t,g}(j) := \overline{Z}_{t,g}^2(j) - \frac{G_t}{t},$

where $\overline{Z}_{t,g}$ is the sample mean of the *g*-th group $\mathcal{Z}_{t,g}$.

Median-of-means type statistic:

$$A_t^{\text{MoM}} := t \cdot \operatorname{median}\left(\sum_{j=1}^p V_{t,1}(j), \sum_{j=1}^p V_{t,2}(j), \dots, \sum_{j=1}^p V_{t,G_t}(j)\right).$$

Test:

$$\phi_{\mathcal{P},\text{dense}} := \mathbb{1}_{\left\{\max_{t\in\mathcal{T}}A_t^{\text{MoM}}/r_t > 1\right\}}.$$

Theorem. Assume $\alpha \geq 2$. Choose $G_t = \min\{t, \Delta\}$ and $r_t = Cp^{(2/\alpha) \vee (1/2)}G_t$, with $\Delta = 8 \log \log(8n)$. Then $\mathcal{R}_{\mathcal{P}}(\rho, \phi_{\mathcal{P}, \text{dense}}) \leq 1/2$ as long as

$$\rho^2 \gtrsim p^{(2/\alpha) \vee (1/2)} \log \log(8n).$$

- When $t \in \mathcal{T} \cap \{t \leq \Delta\}$, MoM simply becomes median.
- When $t \in \mathcal{T} \cap \{t > \Delta\}$, number of groups G_t is at most $\log \log(8n)$.

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- Alternative strategy: for each t ∈ T ∩ {t > Δ̃}, directly apply a robust sparse mean estimator µ̂(·) to

$$\left\{Z_i = (X_i - X_{n+1-i})/\sqrt{2}, \quad i = 1, \dots, t\right\}$$

and use $A_t^{\mathrm{RSM}} := t \| \hat{\mu} \|_2^2$ as the test statistic.

• One example of such estimator $\hat{\mu}(\cdot)$ is given in Prasad et al. (2019):

$$\inf_{\mu \in \mathbb{R}^{p}: \|\mu\|_{0} \leq s} \sup_{u \in \mathcal{N}_{2s}^{1/2}(\mathcal{S}^{p-1})} \left| u^{\top} \mu - 1 \text{DRobust}(\{u^{\top} W_{i}\}_{i=1}^{n}, \eta/(6ep/s)^{s}) \right|,$$

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• One example of such estimator $\hat{\mu}(\cdot)$ is given in Prasad et al. (2019):

 $\inf_{\mu \in \mathbb{R}^p: \|\mu\|_0 \le s} \sup_{u \in \mathcal{N}_{2s}^{1/2}(\mathcal{S}^{p-1})} \left| u^\top \mu - 1 \text{DRobust}(\{u^\top W_i\}_{i=1}^n, \eta/(6ep/s)^s) \right|,$

- 1DRobust: a univariate robust mean est. (e.g. MoM, trimmed mean). - High computational complexity: $|\mathcal{N}_{2s}^{1/2}(\mathcal{S}^{p-1})| \leq (6ep/s)^s$. ▶ We can construct a test $\phi_{\mathcal{P}, \text{sparse}}^{\text{RSM}}$ (non-robust when $t \in \mathcal{T} \cap \{t \leq \tilde{\Delta}\}$) that satisfies $\mathcal{R}_{\mathcal{P}}(\rho, \phi_{\mathcal{P}, \text{sparse}}^{\text{RSM}}) \leq 1/2$ as long as

$$\rho^2 \gtrsim s(p/s)^{2/\alpha} + \log\log(8n).$$

Minimax optimal!

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► To overcome the computation issue, we only use this test when $p \leq \log^{\alpha-2}(\log(8n))$, and use MoM + thresholding + sample splitting otherwise. Best of both worlds!

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Summary

- Quantify the costs of heavy-tailedness on the fundamental difficulty of change point testing problems for high-dimensional data.
- Under $\mathcal{G}_{\alpha,K}$, a CUSUM-type test achieves minimax testing rate up to $\sqrt{\log \log(8n)}$.
- Under $\mathcal{P}_{\alpha,K}$, a median-of-means-type test achieves near-optimal testing rate in both dense and sparse regimes.
- In the sparse regime, a computationally efficient procedure can achieve exact optimality.
- Phase transition at $\alpha = 4$ for $\mathcal{P}_{\alpha,K}$ no sparse regime when $2 \le \alpha \le 4$.

Reference

Li, M.*, Chen, Y.*, Wang, T. and Yu, Y. (2023) Robust mean change point testing in high-dimensional data with heavy tails. *arXiv preprint*, arXiv: 2305.18987.

Thank you!

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