

Iterated Belief Revision and Improvement Operators

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- A Propositional Language \mathcal{L} build from:
 - A set of propositional variables $\mathcal{P} = \{a, b, c, \dots\}$
 - And a set of logical connectives $\wedge, \vee, \neg, \rightarrow, \dots$
- An interpretation $\omega \in \mathcal{W}$ is a function from \mathcal{P} to $\{0, 1\}$ (False/True)
- A model of a formula $\varphi \in \mathcal{L}$ is an interpretation that makes the formula true:

$$\text{mod}(\varphi) = \{\omega \in \mathcal{W} \mid \omega \models \varphi\}$$

- A belief base is a set of propositional formulas that represent the beliefs of the agent
- In the following we will consider that a belief base is a unique propositional formula (the conjunction of the set)

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 - Non-prioritized revision [Theoria 98]

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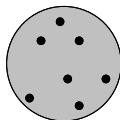
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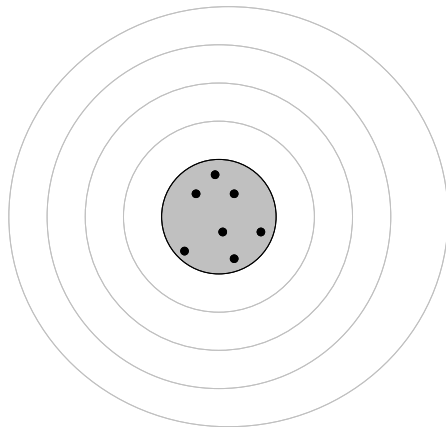
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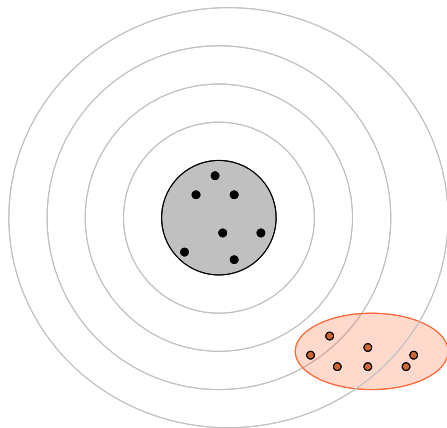


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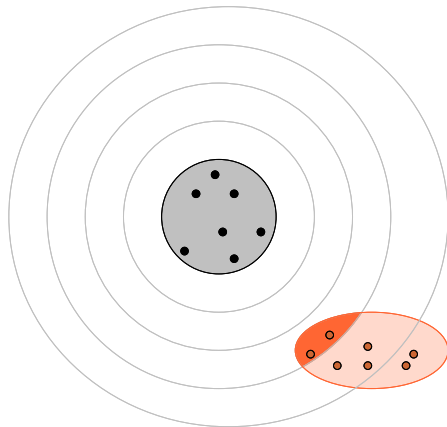


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Revision [Alchourrón-Gärdenfors-Makinson 85]

- (R1)** $\varphi \star \alpha \vdash \alpha$
- (R2)** If $\varphi \wedge \alpha$ is consistent then $\varphi \star \alpha \equiv \varphi \wedge \alpha$
- (R3)** If α is consistent then $\varphi \star \alpha$ is consistent
- (R4)** If $\varphi_1 \equiv \varphi_2$ and $\alpha_1 \equiv \alpha_2$ then $\varphi_1 \star \alpha_1 \equiv \varphi_2 \star \alpha_2$
- (R5)** $(\varphi \star \alpha) \wedge \beta \vdash \varphi \star (\alpha \wedge \beta)$
- (R6)** If $(\varphi \star \alpha) \wedge \beta$ is consistent then $\varphi \star (\alpha \wedge \beta) \vdash (\varphi \star \alpha) \wedge \beta$

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Theorem (Katsuno-Mendelzon 91)

An operator \star is a revision operator (ie. it satisfies (R1)-(R6)) if and only if there exists a faithful assignment that maps each base φ to a total pre-order \leq_φ such that

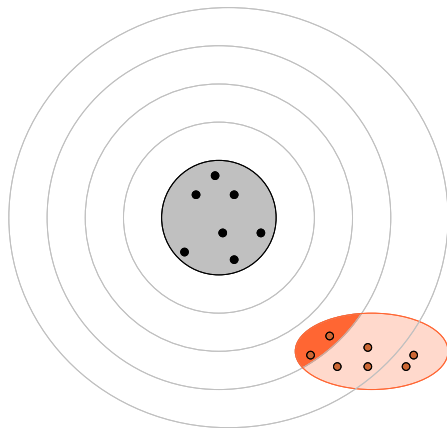
$$\text{mod}(\varphi \star \alpha) = \min(\text{mod}(\alpha), \leq_\varphi).$$

A faithful assignment is a function mapping each base φ to a pre-order \leq_φ over interpretations such that:

- If $\omega \models \varphi$ and $\omega' \models \varphi$, then $\omega \simeq_\varphi \omega'$
- If $\omega \models \varphi$ and $\omega' \not\models \varphi$, then $\omega <_\varphi \omega'$
- If $\varphi \equiv \varphi'$, then $\leq_\varphi = \leq_{\varphi'}$

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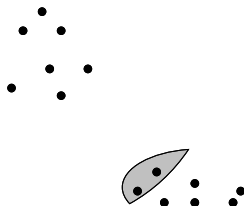
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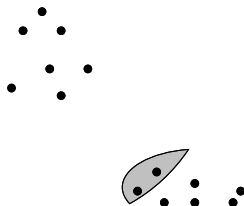


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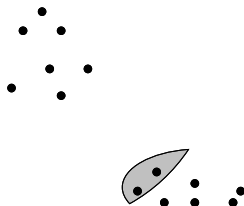
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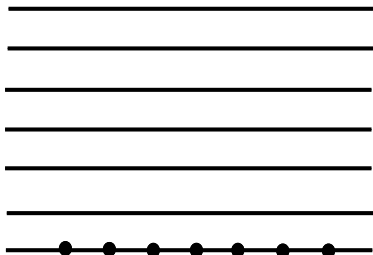
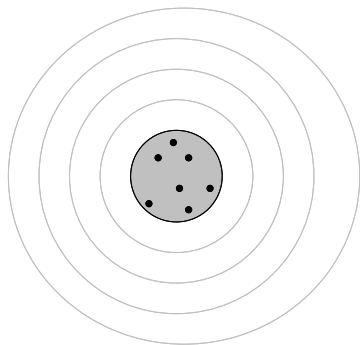
Iterated Belief Revision

- [Darwiche & Pearl 97]
 - An *epistemic state* Ψ encodes the belief base ($B(\Psi)$) + the plausibility pre-order (conditional information)
 - Add constraints on iteration



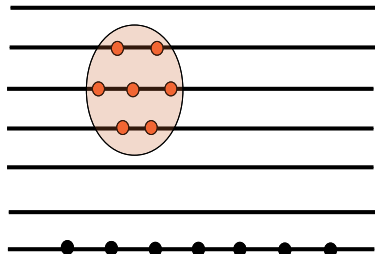
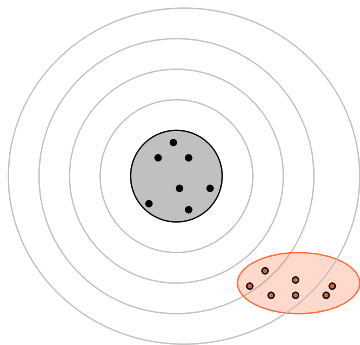
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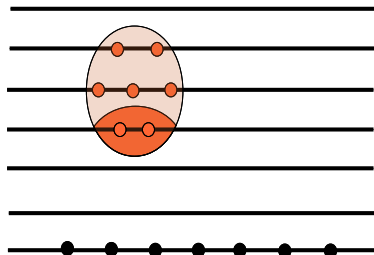
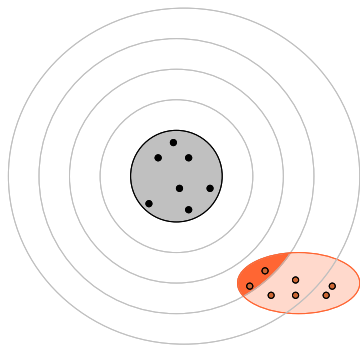
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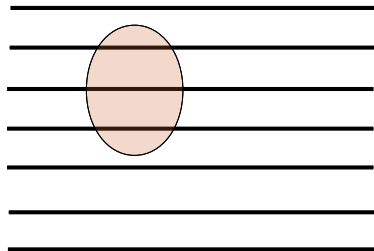
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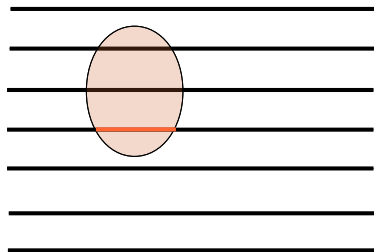
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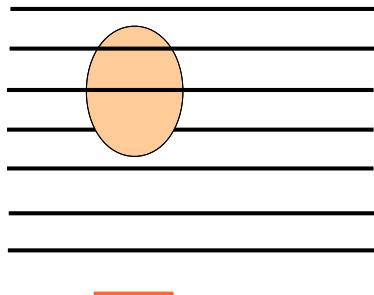
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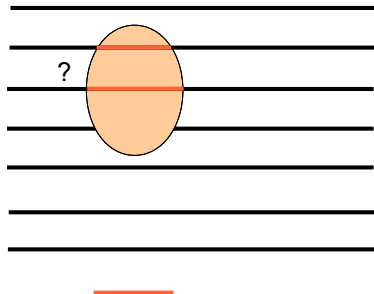


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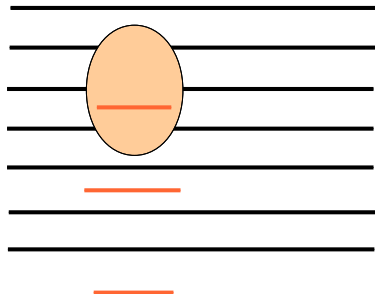


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 - Improve plausibility of models of the new information α

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then $\omega \leq_{\Psi} \omega'$ iff $\omega \leq_{\Psi * \alpha} \omega'$

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Admissible operators:

(CR5) If $\omega \models \alpha$ and $\omega' \models \neg \alpha$,
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- Boutilier's natural revision [Boutilier 96]

Iterated Belief Revision

- [Jin & Thielscher 05,07, Booth & Meyer 06]
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- Boutilier's natural revision [Boutilier 96]
- Nayak's lexicographic operator [Nayak 94, Spohn 88]

Improvement

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 - Increase the plausibility of the new information (CR5)
 - No Primacy of Update

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where n is the first integer such that $\Psi \circ^n \alpha \vdash \alpha$

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Weak Improvement Operators

Definition

An operator \circ is said to be a weak improvement operator if it satisfies **(I1-I6)**:

- (I1)** There exists n such that $\Psi \circ^n \alpha \vdash \alpha$
- (I2)** If $\Psi \wedge \alpha \not\vdash \perp$, then $\Psi \star \alpha \equiv \Psi \wedge \alpha$
- (I3)** If $\alpha \not\vdash \perp$, then $\Psi \circ \alpha \not\vdash \perp$
- (I4)** For any positive integer n if $\alpha_i \equiv \beta_i$ for all $i \leq n$ then
$$\Psi \circ \alpha_1 \circ \dots \circ \alpha_n \equiv \Psi \circ \beta_1 \circ \dots \circ \beta_n$$
- (I5)** $(\Psi \star \alpha) \wedge \beta \vdash \Psi \star (\alpha \wedge \beta)$
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Theorem

A change operator \circ is a weak improvement operator if and only if there exists a strong faithful assignment that maps each epistemic state Ψ to a total pre-order on interpretations \leq_Ψ such that

$$\text{mod}(\Psi \star \alpha) = \min(\text{mod}(\alpha), \leq_\Psi)$$

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Corollary

*If \circ is a weak improvement operator, then \star is an AGM/DP revision operator, i.e. it satisfies (R*1)-(R*6).*

$$\Psi \star (\alpha \vee \beta) = \begin{cases} \Psi \star \alpha \text{ or} \\ \Psi \star \beta \text{ or} \\ \Psi \star \alpha \vee \Psi \star \beta \end{cases}$$

Definition

- We say that α is below β with respect to Ψ , given \circ , denoted $\alpha \prec_{\Psi}^{\circ} \beta$ if and only if $\Psi \star \alpha \vdash \Psi \star (\alpha \vee \beta)$ and $\Psi \star \beta \not\vdash \Psi \star (\alpha \vee \beta)$
- The pair (α, β) is Ψ -consecutive, denoted $\alpha \prec_{\Psi}^{\circ} \beta$ if and only if $\alpha \prec_{\Psi} \beta$ and there is no formula γ such that $\alpha \prec_{\Psi} \gamma \prec_{\Psi} \beta$.

Towards Improvement Operators

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A weak improvement operator is said to be an improvement operator if it satisfies **(I7 - I11)**:

(I7) If $\alpha \vdash \beta$ then $(\Psi \circ \beta) \star \alpha \equiv \Psi \star \alpha$

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Definition

Let \circ be a weak improvement operator and $\Psi \mapsto \leq_{\Psi}$ its corresponding strong faithful assignment. A gradual assignment satisfies:

(S1) If $\omega, \omega' \in \text{mod}(\alpha)$ then $\omega \leq_{\Psi} \omega' \Leftrightarrow \omega \leq_{\Psi \circ \alpha} \omega'$

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small change

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(S3) If $\omega \in \text{mod}(\alpha)$, $\omega' \in \text{mod}(\neg\alpha)$ then $\omega \leq_{\Psi} \omega' \Rightarrow \omega <_{\Psi \circ \alpha} \omega'$ ([JT07,BM06])

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Representation Theorem

Theorem

A change operator \circ is an improvement operator if and only if there exists a gradual assignment such that

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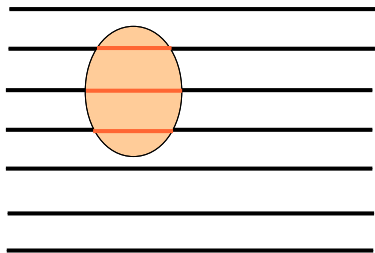
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Proposition

Let \circ be an improvement operator and $\Psi \mapsto \leq_{\Psi}$ its gradual assignment. Then for every formula α , the pre-order $\leq_{\Psi \circ \alpha}$ is completely determined by \leq_{Ψ} and $\text{mod}(\alpha)$.

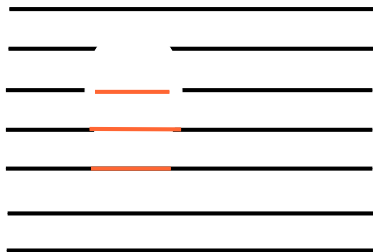
Example

ψ



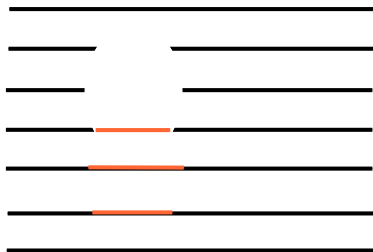
Example

$$\Psi \circ \alpha$$



Example

$$\Psi \circ \alpha \circ \alpha$$



Example

$$\Psi \circ \alpha \circ \alpha \circ \alpha$$

$$\equiv \Psi \dot{-} \neg\alpha$$



Example

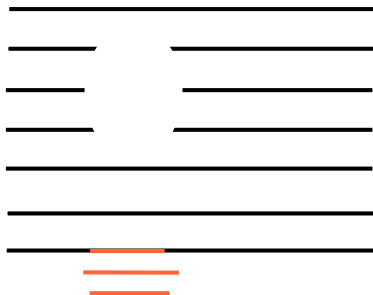
$$\Psi \circ \alpha \circ \alpha \circ \alpha \circ \alpha$$

$$\equiv \Psi \star \alpha$$



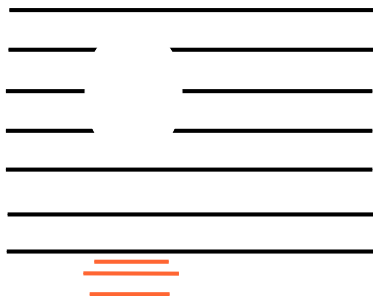
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$$\Psi \circ \alpha \circ \alpha \circ \alpha \circ \alpha \circ \alpha$$



Example

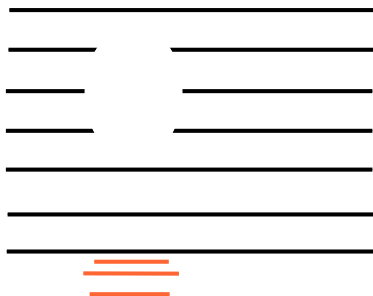
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- ▶ Fix point:
Nayak's lexicographic operator

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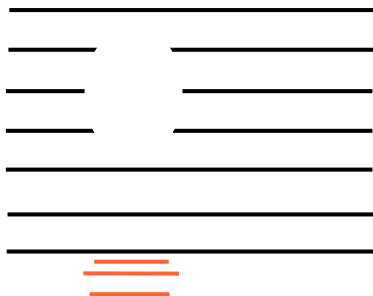
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$\Psi \circ \alpha \circ \alpha \circ \alpha \circ \alpha \circ \alpha \circ \alpha \circ \alpha \circ \alpha$



► Fix point:
Nayak's lexicographic operator

- None of the iterated revision operators has a good long term behaviour
 - Leads to full meet revision
 - ▶ [Lehmann 95]
 - Leads to maxichoice revision
 - ▶ [Darwiche & Pearl 97, Nayak 94, Boutilier 96, Konieczny & Pino Pérez 00, Jin & Thielscher 07, Booth & Meyer 06]

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- None of the iterated revision operators has a good long term behaviour
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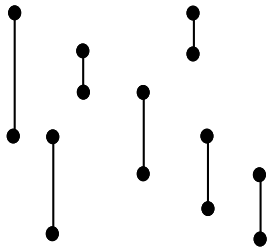
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 - Thanks to the additional input data
- Improvement operators have a good long term behaviour
 - Avoid full meet and maxichoice revision
 - In some iterations new “levels” are created, in some iterations some “levels” are removed.

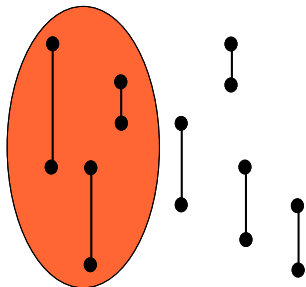
Conclusion

- New family of change operators
 - Increase of the plausibility of the new information
 - No Primacy of Update
- Representation Theorem
- Good long term behaviour
- Others weak improvement operators
- Links with the good day/bad day approach [Booth, Meyer, Wong 06, Booth & Meyer 07]

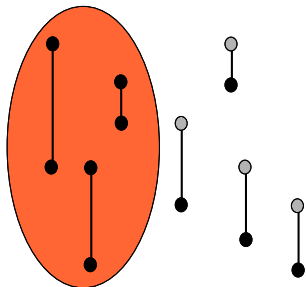
Good Day / Bad Day [Booth, Meyer, Wong 06]



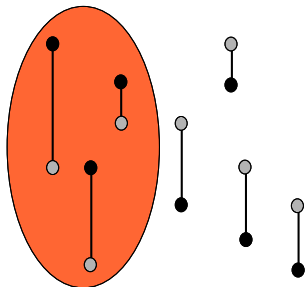
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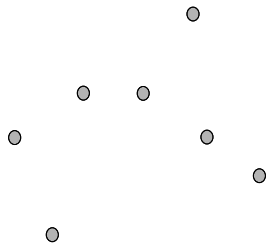


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Good Day / Bad Day [Booth, Meyer, Wong 06]

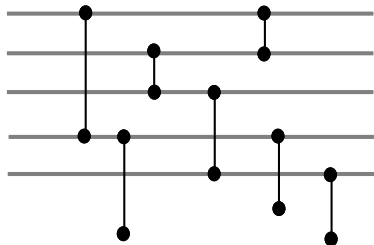




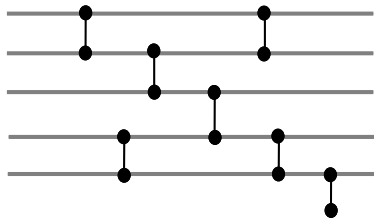
► Iteration ?

- [Booth & Meyer 2007]
- Requires to maintain ($<$, $<<$, $<<<$)

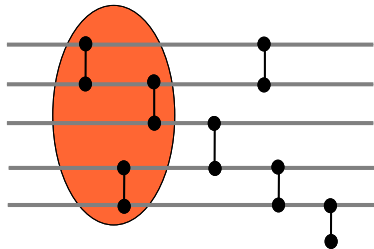
Good Day / Bad Day [Booth, Meyer, Wong 06]



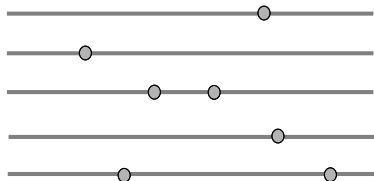
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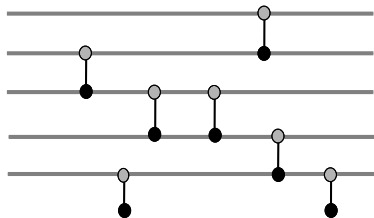
Good Day / Bad Day [Booth, Meyer, Wong 06]



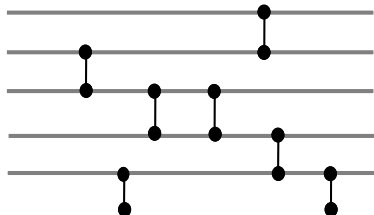
Good Day / Bad Day [Booth, Meyer, Wong 06]



Good Day / Bad Day [Booth, Meyer, Wong 06]



Good Day / Bad Day [Booth, Meyer, Wong 06]



- Improvement operator is a particular case of good day / bad day Revision
 - More general
 - Requires more information

From \leq_Ψ to $\leq_{\Psi \circ \alpha}$

$w \in \text{mod}(\alpha)$	$w' \in \text{mod}(\alpha)$	$w \leq_\Psi w' \Leftrightarrow w \leq_{\Psi \circ \alpha} w'$	(S1)
$w \in \text{mod}(\neg\alpha)$	$w' \in \text{mod}(\neg\alpha)$	$w \leq_\Psi w' \Leftrightarrow w \leq_{\Psi \circ \alpha} w'$	(S2)
$w \in \text{mod}(\alpha)$	$w' \in \text{mod}(\neg\alpha)$	$w <_\Psi w' \Leftrightarrow w <_{\Psi \circ \alpha} w'$ (S3) $w \simeq_\Psi w' \Rightarrow w <_{\Psi \circ \alpha} w'$ (S3) $w' \ll_\Psi w \Rightarrow w \simeq_{\Psi \circ \alpha} w'$ (S4) & (S5) $w' <_\Psi w \wedge w' \not\ll_\Psi w \Rightarrow w <_{\Psi \circ \alpha} w'$ (Lemma)	

Irrelevance of Syntax

Darwiche and Pearl (R*4):

◀ Back

$$\text{If } \alpha \equiv \beta, \text{ then } B(\Psi \circ \alpha) \equiv B(\Psi \circ \beta) \quad (1)$$

Booth and Meyer [06]:

$$\text{If } (\alpha \equiv \beta \ \& \ \gamma \equiv \theta), \text{ then } B(\Psi \circ \alpha \circ \gamma) \equiv B(\Psi \circ \beta \circ \theta) \quad (2)$$

(I4):

For any positive integer n if $\alpha_i \equiv \beta_i$ for all $i \leq n$ then (3)

$$B(\Psi \circ \alpha_1 \circ \dots \circ \alpha_n) \equiv B(\Psi \circ \beta_1 \circ \dots \circ \beta_n)$$

Example (Modification of [Booth & Meyer 06])

Let $\varphi_1 = p \vee \neg p$ and $\varphi_2 = q \vee \neg q$. Let Ψ such that $\Psi \circ \varphi_1 \neq \Psi \circ \varphi_2$. Note that this is compatible with RAGM plus (2), because the only constraint imposed by (2) is that $\leq_{\Psi \circ \varphi_1} = \leq_{\Psi \circ \varphi_2}$ but not that $\Psi \circ \varphi_1 = \Psi \circ \varphi_2$. Moreover we can take $B(\Psi \circ \varphi_1) = p \wedge q = B(\Psi \circ \varphi_2)$. Let Ψ_1, Ψ_2 be two epistemic states such that $B(\Psi_1) = p = B(\Psi_2)$, $00 <_{\Psi_1} 01$ and $01 <_{\Psi_2} 00$. Now, it is compatible with RAGM plus (2) to put $\leq_{\Psi \circ \varphi_1 \circ p} = \leq_{\Psi_1}$ and $\leq_{\Psi \circ \varphi_2 \circ \neg p} = \leq_{\Psi_2}$. But then, by the representation, we have $B(\Psi \circ \varphi_1 \circ p \circ \neg p) = \neg p \wedge \neg q$ and $B(\Psi \circ \varphi_2 \circ \neg p \circ \neg p) = \neg p \wedge q$, what clearly is a counter-example to (I4).

Plausibility pre-order

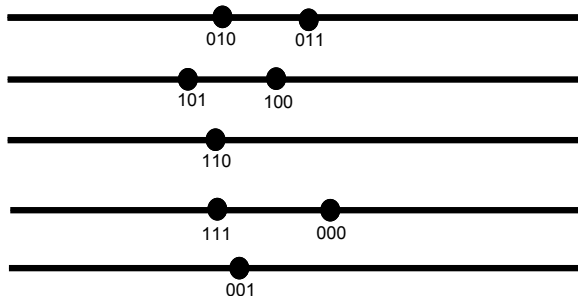
Consider $\mathcal{P} = \{b, f, r\}$

b : the animal is a bird

f : the animal flies

r : the animal is red

$$B(\Psi) = \neg b \wedge \neg f \wedge r$$



Plausibility pre-order

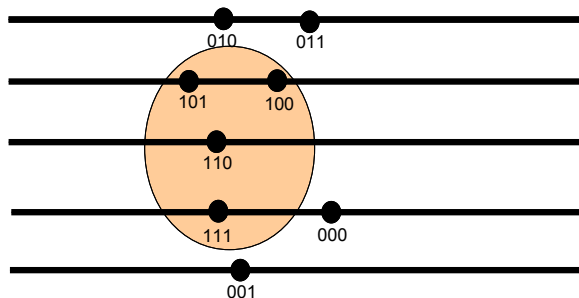
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Plausibility pre-order

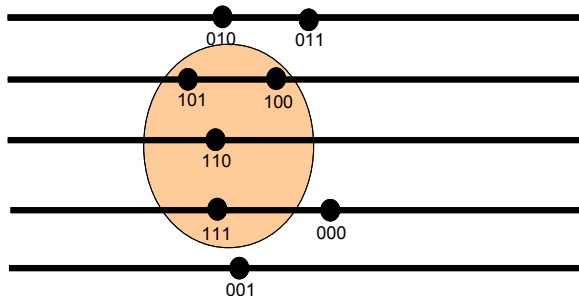
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- $\Psi \star b \vdash f$

Plausibility pre-order

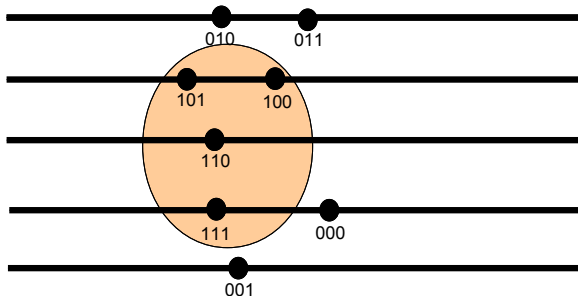
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- $\Psi \star b \vdash f$
- Ramsey test:
“First add the antecedent (hypothetically) to your stock of beliefs ; second make whatever adjustments are required to maintain consistency (without modifying the hypothetical belief in the antecedent); finally, consider whether or not the consequent is true.”

Plausibility pre-order

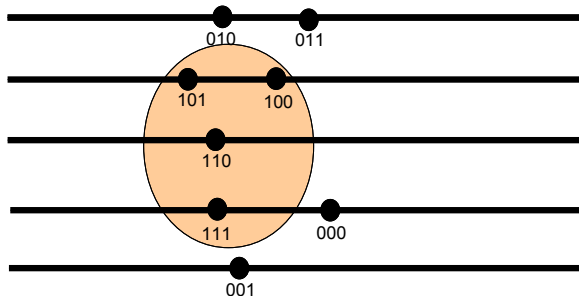
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$$\bullet \quad \Psi \star b \vdash f \quad \Longleftrightarrow \quad f \mid b \in \Psi$$

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- Conditional information

Plausibility pre-order

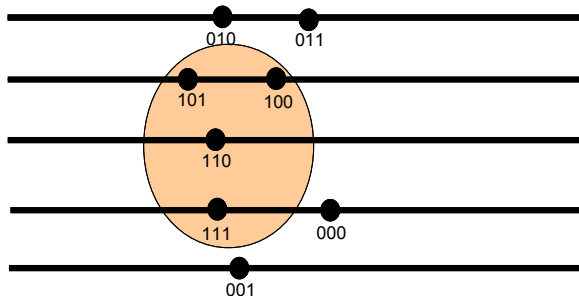
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- Conditional information
- Part of the agent's epistemic state